

## REMARKS

Claims 1-21 are pending in the present application. Claims 1, 6, 10, 15, and 19 are the independent claims. In the Official Action, dated February 9, 2005, claims 1-21 were rejected under § 102(e) as being allegedly anticipated by U.S. Patent No. 6,145,089 (Le et al.).

Initially, the Applicants want to thank the Examiner for his insightful comments. However, Applicants believe that claims 1-21, pursuant to the remarks made below, are patentable over Le et al.

### *Summary of the Invention*

This invention relates to server failover where data is partitioned among a number of servers. For constantly changing data, the data is more typically partitioned across a number of different web servers so that each web server only handles a percentage of all the data. For example, in a two data server scenario, the data of a first type is stored on a first server, and data of a second type is stored on a second server. Thus, the data is said to be partitioned over the first and second servers. The first server services client requests for data of the first type, whereas the second server services client requests for data of the second type. (Application, ¶ 6, lines 1-8).

In a failover scenario, when one of the these two servers is offline, the other server handles its client requests. So that when the first server is offline, the second server becomes the failover server, processing client requests for data usually cached by the first server. And vice versa. The failover server obtains the requested data from the database, temporarily caches the data, and returns the data to the requestor client. When the offline server is back online, and the failover server is notified of this, preferably the failover server then deletes the data it temporarily has cached. Needless to say, this failover procedure applies not only to a two server scenario, but to a plurality of servers so that servers can be bunched into a failover group. (Application, ¶ 8, lines 1-9, and ¶ 45 line 4-5).

In a preferred embodiment, when a server receives a client request, it first determines whether the request is for data of the type normally processed by the server. If it is, the server

processes the request, returning the requested data back to the requestor client. If the data is not normally of the type processed by the server, the server determines whether the correct server to handle data of the type requested has been marked offline. If the correct server has not been marked offline, the server attempts to contact the correct server itself. If successful, the server passes the request to the correct server, which processes the request. If unsuccessful, then the server processes the request itself, querying the database for the requested data where necessary. (Application, ¶ 10, lines 1-10).

*Le et al.*

Le et al. deals with a method and apparatus for a server fail-over system, where a plurality of servers provide a plurality of services. The plurality of servers includes a first server for providing a first service, the system further including a client for consuming the plurality of services, including the first service. If the first server fails to provide the first service, the first service fails over to a second server of the plurality of servers, the second server of the plurality of servers being the highest priority server for providing the first service in the event of failure of the first server. (Abstract).

Specifically, according to Figure 1B, when a failed server C is no longer accessible to the client connected to the network supported by server A, server B, and previously server C, servers A and B continue to support the services previously provided by server C. Thus, for example, while server A continues to support an intranet web server and an NFS server as it did before, after the failover it additionally supports an internet web server previously supported by server C. Selecting whether server A or server B will support the internet web server depends on the priorities of server A and server B and is determined by a nomination and an election process. (Col. 2, lines 37-63 and col. 3, lines 21-23).

For example, the role manager (RM) nominates its server, server A, for each service which server A can provide. The RM has a list of services which server A can provide. When the nomination state is completed, and all servers have converged to a single nominee server, the RM continues to proceed to the election process. The election state response depends on the results of the nomination state. There are several possible results from the nomination state: (1)

time-out on the nomination process, (2) another server already has provided the service, (3) the RM is unable to provide the service, or (4) server A is selected as the server which will provide the service. This process is repeated when the next failover situation comes up. (col. 6, lines 16-63, and Figs. 6 and 9).

***Rejection of Claims 1-21 under 35 U.S.C. § 102(e)***

Claim 1, representative of the other independent claims, 6, 10, 15, and 19, reads as follows:

A system comprising:  
a plurality of servers organized into one or more failover groups and over which ***data is partitioned***, each server usually processing client requests for ***data of a respective type*** and processing the client requests for ***data other than the respective type*** for other of the plurality of servers within a same failover group when the other of the plurality of servers within the same failover group are offline....

(emphasis added). The following passage in the specification further explains partitioning in a two server scenario:

In a two data server scenario, ***data of a first type*** is stored on a first server, and ***data of a second type*** is stored on a second server. It is said that the data of both types is partitioned over the first and the second servers.

(Application, p.2, ¶ 6) (emphasis added). It is the ability of a server, in a failover scenario, to handle a ***data of a second type*** when in its usual operation it handles ***data of a first type*** that distinguishes the claimed invention over Le et al. (*see* claim 1 above in conjunction with the accompanying passage).

Le et al. concerns the failover of plurality of ***services*** (not different ***types of data***). Thus, if a first server fails, a second server will provide a service previously provided by the first service (Fig. 1, and col. 2, ll. 22-63). Le et al. says nothing about the second server having the ability to take over different ***types of data*** that were normally provided by the first server. Put another way, the second server could be providing services of data type X and then take over *different services* also of data type X from the failed first server – and still not be able to handle *different*

*types of data* (for example, data of type Y). In short, Le et al. does not teach the failover of different *data types* (but merely that of plurality of *services*).

Thus, Le et al. fails to disclose: “server[s] usually processing client requests for *data of a respective type* and processing the client requests for *data other than the respective type* for other of the plurality of servers within a same failover group...” (claim 1)(emphasis added).

Likewise, Le et al. fails to disclose similar limitations present in the other independent claims 6, 10, 15, and 19: “a plurality of servers organized into one or more failover groups, each server usually processing client *requests of a respective type* and processing the client requests *other than the respective type* for other of the plurality of servers within a same failover group when the other of the plurality of servers within the same failover group are offline” (claim 6) (emphasis added); “sending the request to the failover server, capable of processing *requests for partitioned data of a respective type* and *partitioned data other than its respective type*” (claim 10)(emphasis added); “determining whether the *request is of a type* usually processed by the server; in response to determining that the request is of the type usually processed by the server, processing the request; in response to determining that the *request is not of the type* usually processed by the server, determining whether a second server that usually processes the type of the request is indicated as offline” (claim 15)(emphasis added); “notifying a plurality of servers, capable of processing *requests for partitioned data of a respective type* and partitioned *data other than its respective type*, other than the server marked as offline that the server is offline” (claim 19)(emphasis added).

Accordingly, it is respectfully submitted that claim 1 and the other independent claims, 6, 10, 15, and 19, are patentable over Le et al. Since dependent claims 2-5, 7-9, 11-14, 16-18, and 20-21 depend either directly or indirectly from independent claims 1, 6, 10, 15, and 19, respectively, they are believed allowable for the same reasons. Withdrawal of the rejection under § 102(e) is therefore earnestly solicited.

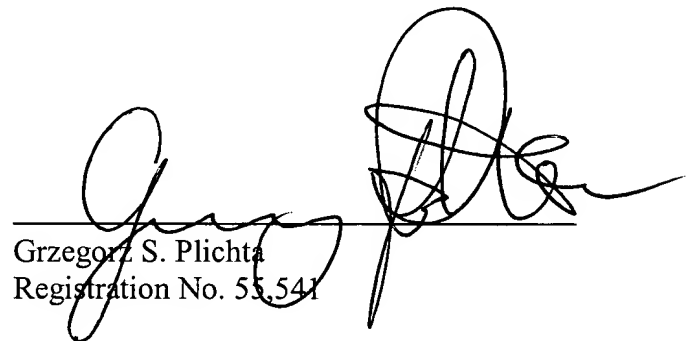
**DOCKET NO.:** MSFT-0593/158452.01  
**Application No.:** 09/681,309  
**Office Action Dated:** February 9, 2005

**PATENT**  
**REPLY FILED UNDER EXPEDITED**  
**PROCEDURE PURSUANT TO**  
**37 CFR § 1.116**

### **CONCLUSION**

Applicants believe that the present Amendment is responsive to each of the points raised by the Examiner in the Official action, and submits that Claims 1-21 of the application are in condition for allowance. Favorable consideration and passage to issue of the application at the Examiner's earliest convenience is earnestly solicited.

Date: March 14, 2005



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